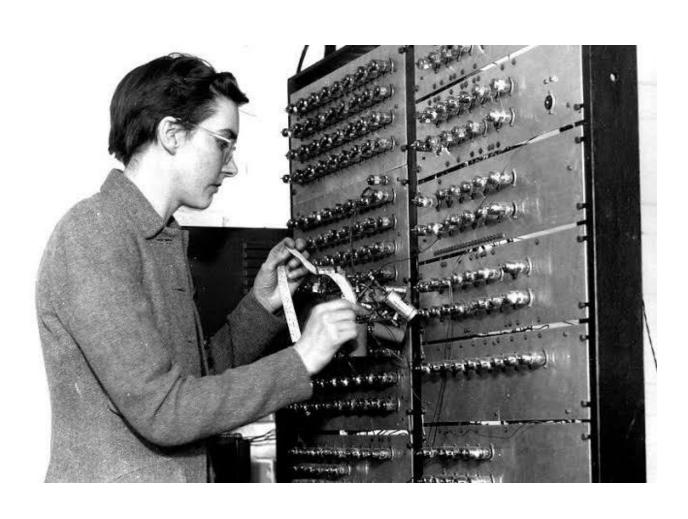
# CSCI 210: Computer Architecture Lecture 9: Logical Operations

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## CS History: Kathleen Britton



- Applied mathematician and computer scientist
- Wrote the first assembly language and assembler in 1947
- Collaborated with Andrew Booth to develop three early computers: the ARC (Automatic Relay Calculator), SEC (Simple Electronic Computer), and APE(X)C
- Later worked with neural nets

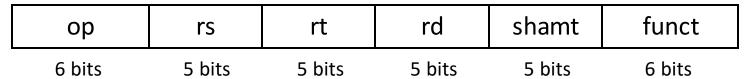
## **Logical Operations**

Instructions for bitwise manipulation

Operation	С	Java	MIPS
Shift left	<<	<<	sll
Shift right	>>	>>>	srl
Bitwise AND	&	&	and, andi
Bitwise OR			or, ori
Bitwise NOT	~	~	nor

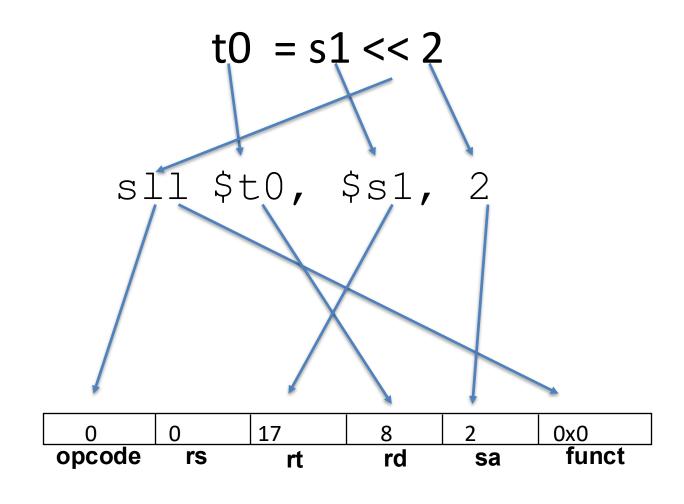
Useful for extracting and inserting groups of bits in a word

## **Shift Operations**



- shamt: how many positions to shift
- Shift left logical
  - Shift left and fill with 0 bits
  - sll by n bits multiplies by 2<sup>n</sup>
- Shift right logical
  - Shift right and fill with 0 bits
  - srl by n bits divides by 2<sup>n</sup> (unsigned only)

### MIPS shift instructions



## Shift left logical

- 0110 1001 << 2 in 8 bits
  - Most significant 2 bits are dropped
  - 2 Os are added to become the least significant bits
  - Result: 01 1010 0100 => 1010 0100

## Shift right logical

- 1010 1001 >>> 3 in 8 bits
  - Least significant 3 bits are dropped
  - 3 Os are added to become the most significant bits
  - Result: 0001 0101 <del>001</del> => 0001 0101

## Shift right arithmetic

- sra rd, rt, shamt
  - Shift right and copy the sign bit
- 1010 1001 >> 3 in 8 bits
  - Least significant 3 bits are dropped
  - 3 1s are added because the MSB is 1 to become the most significant bits
  - Result: 1111 0101 001 => 1111 0101

# A new op HEXSHIFTRIGHT shifts hex numbers right by a digit. HEXSHIFTRIGHT *i* times is equivalent to

A. Dividing by i

B. Dividing by  $2^i$ 

C. Dividing by 16<sup>i</sup>

D. Multiplying by 16<sup>i</sup>

## Remember Boolean Operations?

• and, or, not . . .

Now we'll apply them to bits!

Just think of 1 as True, and 0 as False

## And Truth Table

	0	1
0	0	0
1	0	1

## **AND Operations**

- Useful to mask bits in a word
  - Select some bits, clear others to 0

```
and $t0, $t1, $t2
```

```
$t2 | 0000 0000 0000 0000 1101 1100 0000
```

```
$t0 | 0000 0000 0000 0000 1100 0000 0000
```

## AND identities (for a single bit)

• x & 0 =

• x & 1 =

01101001

& 11000111

A. 00010000

B. 01000001

C. 10101110

D. 11101111

## If we want to zero out bits\* 3 – 0 in a byte we should AND with

A. 00000000

\*MSB (bit 7) is on the left, LSB (bit 0) is on the right

B. 00001111

C. 11110000

D. 11111111

One way to represent colors is to specify the amount of red, green, and blue (RGB) that makes up the color. The three color "channels" are often packed into a 32-bit integer as follows: 0000000 rrrrrrr gggggggg bbbbbbbb where each color is represented in 8 bits. If color is a 32-bit RGB color, which high-level expression extracts the green channel as an 8-bit value in the range 0–255?

A.green = (color >> 8) & 
$$0xFF$$
  
B.green = (color >> 16) &  $0xFF$   
C.green = (color &  $0xFF$ ) >> 8  
D.green = (color &  $0xFF00$ ) >> 16

E. More than one of the above (which?)

Assume color is stored in \$t0 and green should be extracted to \$t1. Which sequence of MIPS instructions corresponds to green = (color >> 8) & 0xFF

```
A.sll $t1, $t0, 8
  andi $t1, $t1, 0xFF
B.srl $t1, $t0, 8
  andi $t1, $t1, 0xFF
C.andi $t1, $t0, 0xFF
  sll $t1, $t1, 8
D.andi $t1, $t0, 0xFF
  srl $t1, $t1, 8
```

## Or Truth Table

	0	1
0	0	1
1	1	1

## **OR Operations**

- Useful to set bits in a word
  - Set some bits to 1, leave others unchanged

```
or $t0, $t1, $t2
```

```
$t2 | 0000 0000 0000 0000 1101 1100 0000
```

```
$t1 | 0000 0000 0000 0001 1100 0000 0000
```

\$t0 | 0000 0000 0000 00011 1101 1100 0000

## OR Identities (for a single bit)

01101001 11000111

A.00010000

B.0100001

C.10101110

D.11101111

#### Recall RGB:

0000000 rrrrrrr ggggggg bbbbbbbb

If r, g, and b are values in the range 0–255, how can we construct the RGB value, c, whose channels are r, g, and b?

A.c = 
$$(r << 16) \mid (g << 8) \mid b$$
  
B.c =  $(r << 24) \mid (g << 16) \mid (b << 8)$   
C.c =  $(r >> 8) \mid (g >> 16) \mid (b >> 24)$   
D.c =  $(r >> 16) \mid (g >> 24) \mid (b >> 32)$ 

E. More than one of the above (which?)

## Nor Truth Table

	0	1
0	1	0
1	0	0

### **NOR Operations**

MIPS has NOR 3-operand instruction

```
-a NOR b = NOT (a OR b)
```

nor \$t0, \$t1, \$t2

```
$t2 | 0000 0000 0000 0000 1101 1100 0000
```

```
$t1 | 0000 0000 0000 0001 1100 0000 0000
```

```
$t0 | 1111 1111 1111 1110 0010 0011 1111
```

## 01101001 NOR 11000111

A. 00010000

B. 01000001

C. 10101110

D. 11101111

## **NOT** operations

- Inverts all the bits in a word
  - Change 0 to 1, and 1 to 0

## MIPs does not need a NOT operation because we can use \_\_\_\_ for NOT \$t1, \$t2

A. nor \$t1, \$t2, \$zero

B. nor \$t1, \$t2, \$t3, where all bits in \$t3 are set to 1

C.nori \$t1, \$t2, 0b1111111\_1111111111, where nori is Norlmmediate

- D. It does require a NOT operation
- E. None of the above are correct

## **XOR Truth Table**

	0	1
0	0	1
1	1	0

## **XOR Operations**

- Exclusive OR (written x ⊕ y or x ^ y)
  - Set bits to one only if they are not the same

```
xor $t0, $t1, $t2
```

```
$t2 | 0000 0000 0000 0000 1101 1100 0000
```

\$t0 | 0000 0000 0000 0001 0001 1100 0000

### 01101001 XOR 11000111

A. 00010000

B. 01000001

C. 10101110

D. 11101111

## XOR Identities (for a single bit)

• x XOR 0 =

• x XOR 1 =

## 10 & 7

A. 0

B. 2

C. 7

D. 10

E. None of the above

# Set bit 4 in byte x to 1, leaving the rest of the bits unchanged

A. x = x AND 00010000

B. x = x AND 11101111

C. x = x OR 00010000

D. x = x NOR 11101111

### Invert bits 2–0 of x

A. x = x AND 00000111

B. x = x OR 00000111

C. x = x NOR 00000111

D. x = x XOR 00000111

## Find the ones' complement of x (in 8 bits)

A. x XOR 00000000

B. x XOR 11111111

C. x XOR 11111110

D. x OR 11111111

## Reading

• Next lecture: Branching instructions